

An Introduction to Genetic Algorithms

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OUTLINE

- (Historical background
- (Working Principles of GA
- (GA operators
- (Concept of schemata in GA
- (GA v/s traditional methods
- (Drawbacks of GA
- (GA for constrained optimization
- (Modern developments

Historical Background

(Mechanics of natural genetics and evolution

- u Survival of the fittest
- u Evolution of the species

(John Holland, University of Michigan

- s *Adaptation in natural and artificial systems*, 1975

(David Goldberg, University of Illinois

- s *Genetic Algorithms in Search, Optimization and Machine Learning*, 1989

(ICGA, 1985, 87, 89,

Working Principles of GA

(Unconstrained Optimization

- u Maximize $f(\mathbf{x})$, $\mathbf{x} = x_i, i = 1, n, U_i \leq x_i \leq L_i, i = 1, n$

(Coding of \mathbf{x} in string structures

- u Not absolutely essential (Real GA)
- u Usually Binary Coding
 - s each x_i coded in k sub-strings of length $l_i, k = 2^{\{l_i\}}$
 - s sub-string element $s_k(m) \in (0,1), m = 0, l_i-1$
 - $s_k = s_{l_i-1} s_{l_i-2} \dots s_2 s_1 s_0$
 - s Discrete values of x_i mapped
 - $x_i = L_i + DV(s_i) * (U_i - L_i) / (2^{\{l_i\}} - 1)$, where
 $DV(s_k) = \text{Decoded value of } (s_k) = \sum 2^m \cdot s_k(m), m = 0, l_i-1$

Example of Binary Coding

u Decoding Procedure

s $i=1, l_1 = 4$

s $k = \text{total number of discrete values of } x_1 = 2^{\{l_1\}} = 2^4 = 16$

s string $s_k = s_3s_2s_1s_0$

s $s_k(m) \in (0,1), m = 0, 3$

s $DV(s_k) = \sum 2^m \cdot s_k(m), m = 0, 3$

u Accuracy

s $\approx (U_i - L_i) / 2^{\{l_i\}}$

s Exponential increase with string length

| k | s_k | DECODING OF s_k | $x_1(k)$ |
|----|-------|--|----------|
| 1 | 0000 | $(0) \cdot 2^3 + (0)2^2 + (0) \cdot 2^1 + (0) \cdot 2^0$ | 0 |
| 2 | 0001 | $(0) \cdot 2^3 + (0)2^2 + (0) \cdot 2^1 + (1) \cdot 2^0$ | 1 |
| 3 | 0010 | $(0) \cdot 2^3 + (0)2^2 + (1) \cdot 2^1 + (0) \cdot 2^0$ | 2 |
| 4 | 0011 | $(0) \cdot 2^3 + (0)2^2 + (1) \cdot 2^1 + (1) \cdot 2^0$ | 3 |
| 5 | 0100 | $(0) \cdot 2^3 + (1)2^2 + (0) \cdot 2^1 + (0) \cdot 2^0$ | 4 |
| 6 | 0101 | $(0) \cdot 2^3 + (1)2^2 + (0) \cdot 2^1 + (1) \cdot 2^0$ | 5 |
| 7 | 0110 | $(0) \cdot 2^3 + (1)2^2 + (1) \cdot 2^1 + (0) \cdot 2^0$ | 6 |
| 8 | 0111 | $(0) \cdot 2^3 + (1)2^2 + (1) \cdot 2^1 + (1) \cdot 2^0$ | 7 |
| 9 | 1000 | $(1) \cdot 2^3 + (0)2^2 + (0) \cdot 2^1 + (0) \cdot 2^0$ | 8 |
| 10 | 1001 | $(1) \cdot 2^3 + (0)2^2 + (0) \cdot 2^1 + (1) \cdot 2^0$ | 9 |
| 11 | 1010 | $(1) \cdot 2^3 + (0)2^2 + (1) \cdot 2^1 + (0) \cdot 2^0$ | 10 |
| 12 | 1011 | $(1) \cdot 2^3 + (0)2^2 + (1) \cdot 2^1 + (1) \cdot 2^0$ | 11 |
| 13 | 1100 | $(1) \cdot 2^3 + (1)2^2 + (0) \cdot 2^1 + (0) \cdot 2^0$ | 12 |
| 14 | 1101 | $(1) \cdot 2^3 + (1)2^2 + (0) \cdot 2^1 + (1) \cdot 2^0$ | 13 |
| 15 | 1110 | $(1) \cdot 2^3 + (1)2^2 + (1) \cdot 2^1 + (0) \cdot 2^0$ | 14 |
| 16 | 1111 | $(1) \cdot 2^3 + (1)2^2 + (1) \cdot 2^1 + (1) \cdot 2^0$ | 15 |

Working Principles of GA

(Fitness Function $F(\mathbf{x})$

- u Directly related to $f(\mathbf{x})$
 - s $F(\mathbf{x}) = f(\mathbf{x})$ for maximization
 - s $F(\mathbf{x}) = 1 / (1 + f(\mathbf{x}))$ for minimization

(Steps in GA

- u Create initial population (randomly) by concatenating sub-strings
- u Evaluate $F(\mathbf{x})$ of each member, identify best member(s)
- u Create next generation, using GA operations
 - s Reproduction
 - s Crossover
 - s Mutation
- u Iterate till convergence

Reproduction

(Insertion of strings with higher $F(\mathbf{x})$ in mating pool

u Strings with low $F(\mathbf{x})$ may also get selected

(Selection strategies

u Proportionate

s probability of selection proportional to $F(\mathbf{x})$

- Roulette-Wheel selection

u Ranking

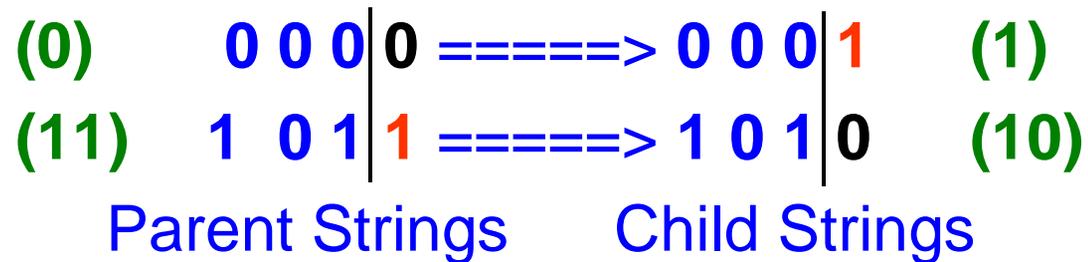
s best of a few strings selected each time

- Tournament selection

Crossover

- u Two strings from mating pool selected (randomly)
- u Location(s) of crossover location determined (randomly)
- u Bits of Strings interchanged between crossover location(s)

(Single point Crossover



(Crossover can be beneficial or detrimental

- u Crossover Probability $p_c \approx 0.7$ to 0.9

Mutation

(Flipping of bits of strings at random

- u local search around current solution
- u maintain genetic diversity of population
- u enable search to climb towards global optimum
- u mutation probability $p_m \approx 0.005$ to 0.015

(Illustration with $p_m = 0.10$

- u Original population

0 0 1 0 , 1 0 0 1 , 0 0 0 1 , 1 1 1 1 , 0 1 1 1 , 1 1 0 1 , 1 1 1 1 , 1 0 0 1
(2) (9) (1) (16) (7) (13) (16) (9)

- u After Mutation

0 0 1 0 , 1 0 0 1 , 0 1 0 1 , 1 1 1 1 , 0 1 1 0 , 1 1 0 1 , 1 1 1 1 , 1 1 0 1
(2) (9) (5) (16) (6) (13) (16) (13)

GA Operators - Summary

(Reproduction

- u Select **good** strings from population (eliminate **bad** strings)

(Crossover

- u Recombine strings and (hopefully) create better strings

(Mutation

- u Once in a while, create even better/worse strings

Concept of Schemata in GA

(Schemata

- s number of strings with similarities at certain string positions
- s total schemata for binary representation = 3^l , l = string length
- s defining length $\delta(H)$ = difference in outermost defining points
- s order $o(H)$ = number of fixed point
- s Example for $H = 1*0^*$
 - $\delta(H) = o(H) = 2$
 - strings = 1101, 1100, 1000, 1001
 - occupies 25% of search space

| | | | |
|------|------|------|------|
| 0000 | 0001 | 0010 | 0011 |
| 0100 | 0101 | 0110 | 0111 |
| 1000 | 1001 | 1010 | 1011 |
| 1100 | 1101 | 1110 | 1111 |

- s population = 16, but total schemata = $3^4 = 81$

Schemata Theorem

$$m(H, t + 1) \geq m(H, t) \frac{F(H)}{F_{avg}} \left[1 - p_c \frac{d(H)}{l-1} - p_m o(H) \right]$$

Where,

$m(H, t)$ = Number of strings of schema H in generation t

$F(H)$ = Average fitness of strings of schema H

F_{avg} = Average fitness of the population

p_c = Probability of crossover

p_m = Probability of mutation

l = String length

$d(H)$ = Defining length of schema H

$o(H)$ = Order of schema H

Building Block Hypothesis

u Building Blocks

- s Schema with low $d(H)$ & $o(H)$, and $F(H) > F_{avg}$
 - Represent different large, good regions in search space
 - Propagate exponentially with generations
 - Combine with each other vigorously
 - Lead to
 - ┆ optimum (near-optimal) solutions

u Implicit Parallelism

- s In each generation
 - n function evaluations
 - n^3 schemata processed
 - without any specific book-keeping !

GA v/s Traditional methods

(Direct Methods

- u Gradients or any auxiliary information not required
- u Can handle discontinuous, Multi-modal functions

(Stochastic formulation

- u Independent of starting point

(Work on coding of variables

- u Domain Discretization
- u General purpose & robust code possible

(Work on a population of points

- u Better chance of converging to global minimum
- u Can identify multiple optimal solutions

Drawbacks of GA

- u Coding scheme should be meaningful & appropriate
 - s premature convergence
- u GA parameters have to be tuned before starting
 - s population size, max. generations, p_c , p_m
 - s crossover type, selection strategy
- u Constraints cannot be implicitly handled
 - s penalty function approach (usually)
- u Large number of function evaluations
 - s Orders of magnitude higher than traditional algorithms

GA for constrained optimization

Maximize $F(\mathbf{x}) = 1/(1+ P(\mathbf{x}))$ or

Minimize $F(\mathbf{x}) = P(\mathbf{x})$, where

$$P(\mathbf{x}) = f(\mathbf{x}) + \sum_{j=1}^J u_j \langle g_j(\mathbf{x}) \rangle^2 + \sum_{k=1}^K v_k \langle h_k(\mathbf{x}) \rangle^2$$

$f(\mathbf{x})$ = objective function

u_j = penalty coefficients for J inequality constraints g_j

v_k = penalty coefficients for K equality constraints h_k

u_j & v_k are usually kept constant during one GA run

Large values can be assigned to u_j & v_k

Other GA operators

u Fitness Function scaling

- s To avoid premature convergence due to
 - dominance by one string
 - presence of several average fit strings
 - | transform $F(x)$ to $S(x) = aF(x) + b$; choose a & b such that
 - | best string has a predefined number of copies
 - | each string with average fit has only one copy

u Multi-point Crossover

- s Single point crossover is biased in favor of right-most bits

u Uniform Crossover

- s A bit at any location chosen either parent with $p = 0.5$
- s for e.g. if 1st and 3rd bits are changed, then

(0) 0 0 0 0 =====> 1 0 1 0 (10)
(11) 1 0 1 1 =====> 0 0 0 1 (1)

Recent development in GAs

- u **Real Coded GA**
 - s work directly on design variables
- u **Simultaneous multiple solutions of multi-modal functions**
 - s divide population using sharing functions
- u **Multi-criteria optimization**
 - s pareto-optimal solutions

- u **Reported applications of GA in literature**
 - s Gas pipeline layout optimization
 - s Job shop scheduling, Routing, Travelling salesperson
 - s Aircraft Design and Structural / Aerodynamic optimization
 - s Gas Turbine blade design
 - s Laminate stacking in Composite structures

Happy computing with GA !!